

Automata and Zappa-Szép products of groups

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Zappa-Szép products

Let $H, K \leq G$.

- If $G = HK$ with $H \cap K = 1$ and $H, K \triangleleft G$ then G is an *internal direct product* of H and K , $G = H \times K$.
- If $G = HK$ with $H \cap K = 1$ and $H \triangleleft G$ then G is an *internal semi-direct product* of H and K , $G = H \rtimes K$.
- If $G = HK$ with $H \cap K = 1$ then G is an *internal Zappa-Szép direct product* of H and K , $G = H \ltimes K$.

Other names: general product, knit product, exact factorization.

there exists unique elements $k' \in K$ and $h' \in H$ such that $kh = \theta_k(h)(k \cdot h)$.

As $G = HK$, if $kh \in G$, $\exists k' \in K, h' \in H$ s.t.
 $kh = h'k'$.

h' depends on h & $k \mapsto \theta_k(h)$

k' depends on h & $h \mapsto k \cdot h$

We therefore have two functions

$$(k, h) \mapsto \theta_k(h) \in H, \quad (k, h) \mapsto k \cdot h \in K$$

called the mutual actions defined by the multiplication.

$$kh = \theta_k(h)(k \cdot h) \quad \psi^k$$

$$\begin{aligned} k h_1 h_2 &= \theta_k(h_1)(k \cdot h_1) h_2 \\ &= \underbrace{\theta_k(h_1) \theta_{k \cdot h_1}(h_2)}_{\parallel \in H} \underbrace{[(k \cdot h_1) \cdot h_2]}_k \\ &\rightarrow \theta_k(h_1 h_2)(k \cdot h_1 h_2) \end{aligned}$$

$$\theta_k(h_1 h_2) = \theta_k(h_1) \theta_{k \cdot h_1}(h_2)$$

$$\& (k \cdot h_1) \cdot h_2 = k \cdot (h_1 h_2) \quad k_1, k_2, h$$

This general idea allows for generalisations to other algebraic structures (see Brin *On the Zappa-Szép product*, Comm. Algebra 2005).

Take these conditions as axioms which two functions must satisfy:

Let H, K be semigroups where $h, h_1, h_2 \in H$ and $k, k_1, k_2 \in K$:

$$\begin{aligned} & \rightarrow \left(\begin{aligned} k \cdot (h_1 h_2) &= (k \cdot h_1) \cdot h_2, \\ \theta_k(h_1 h_2) &= \theta_k(h_1) \theta_{k \cdot h_1}(h_2), \end{aligned} \right) \\ & \left(\begin{aligned} \theta_{k_1 k_2}(h) &= \theta_{k_1}(\theta_{k_2}(h)), \\ (k_1 k_2) \cdot h &= (k_1 \cdot \theta_{k_2}(h))(k_2 \cdot h) \end{aligned} \right) \end{aligned}$$

Then under the multiplication

$$h_1 k_1 h_2 k_2 = h_1 \theta_k(h_2) (k_1 \cdot h_2) k_2$$

$$(h_1, k_1)(h_2, k_2) = (h_1 \theta_{k_1}(h_2), (k_1 \cdot h_2) k_2), \quad \left. \right)$$

the set $H \times K$ defines a semigroup called the *external Zappa–Szép semigroup of H and K* , denoted $H \bowtie K$.

If H and K are monoids and the multiplication satisfies

$$\begin{aligned} & \rightarrow k \cdot 1_H = k \quad \text{'' action} \\ & \theta_{1_K}(h) = h \quad \text{''} \end{aligned} \quad \theta(1_H) = 1_H$$

then $H \bowtie K$ is a ~~monoid~~, called the *external Zappa–Szép monoid of H and K* .

If H and K are groups then the external Zappa–Szép monoid is in fact a group, the *external Zappa–Szép group of H and K* .

1993 Kunze

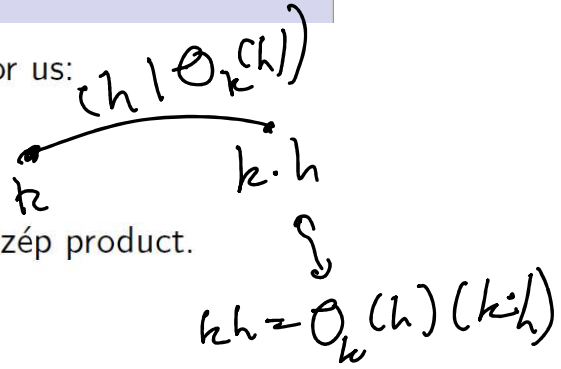
Automata

An automaton $\mathcal{A}_{H \rtimes K}$ is a labeled, directed graph. For us:

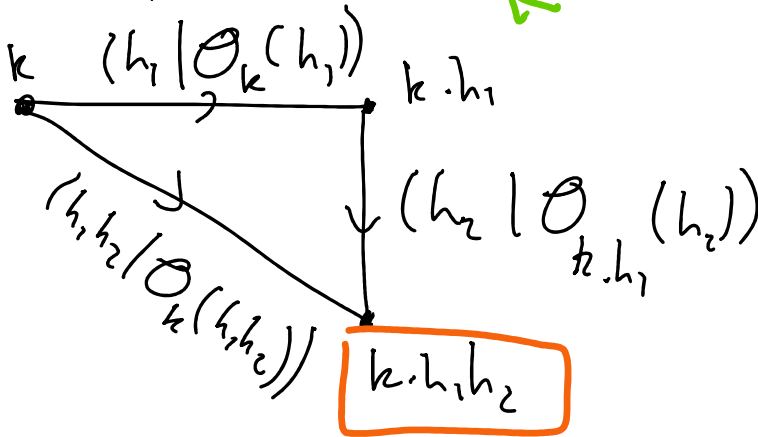
- vertices from (a subset of) K ,
- edge labels from (a subset of) $H \times H$.

Want $\mathcal{A}_{H \rtimes K}$ to encode the structure of the Zappa–Szép product.

$$\begin{aligned} \forall k \cdot (h_1 h_2) &= (k \cdot h_1) \cdot h_2, \\ \theta_k(h_1 h_2) &= \theta_k(h_1) \theta_{k \cdot h_1}(h_2), \\ \theta_{k_1 k_2}(h) &= \theta_{k_1}(\theta_{k_2}(h)), \\ (k_1 k_2) \cdot h &= (k_1 \cdot \theta_{k_2}(h))(k_2 \cdot h) \end{aligned}$$

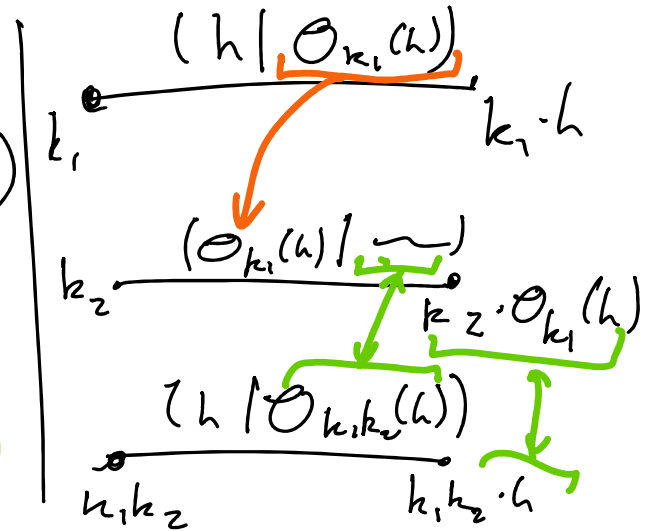


Serial processing \leftarrow



$$\theta_k(h_1) \theta_{k \cdot h_1}(h_2) = \theta_k(h_1, h_2)$$

Sequential processing



Presentations

If H and K are groups then $H \rtimes K$ has relative presentation

$$\langle H, K \mid T_{H \rtimes K} \rangle$$

where $T_{H \rtimes K}$ consists of elements of the form $kh = \theta_k(h)(k \cdot h)$.

Question

Given H, K and $T_{H \bowtie K}$ as above, under what conditions is

$$\langle H, K \mid H, K, T_{H \bowtie K} \rangle$$

a Zappa-Szép semigroup/monoid/group $H \bowtie K$?

$$\langle \overset{H}{a, b}, \overset{K}{x, y} \mid ax = ya, ay = ya, \dots \rangle \rightarrow c = y$$

Plan of action:

- 1 Fix an alphabet X .
- 2 Fix a group H .
- 3 Define $T_{H \bowtie K}$ via an automaton $\mathcal{A}_{(H, X)}$.
- 4 Placing restrictions on $\mathcal{A}_{(H, X)}$ gives:
 - a semigroup $H \bowtie X^+$.
 - a monoid $H \bowtie X^*$.
 - a group $H \bowtie F(X)$.

$$\begin{aligned} K &= X^+ \\ K &= X^* \\ K &= F(X) \end{aligned}$$

Theorem

Suppose $\mathcal{A}_{(H, X)}$ satisfies the serial processing condition. Then

$$\text{Sem}[H, X \mid \underline{XH = \theta_X(h)(X \cdot h)} \ (h \in H, x \in X)] \stackrel{P}{=} (H \bowtie \theta_X(h))$$

is a Zappa-Szép semigroup of H and X^+ .

①

Proof.

The rewriting system defined by $xg \rightarrow \theta_x(g) (xc \cdot g) \quad x \in X$
 is complete $\rightarrow \forall w \in X^+$
 $\dots \rightarrow w = a'w'$

FS complete $\rightarrow \forall w \in X^*$
 ($\exists! g' \in H \ \& \ w' \in X^* \text{ s.t. } w g = g' w'$)

$$w g = w_1 x_1 g = w_1 \theta_{x_1} (g) (x_1 \cdot g)$$

② So the map $\rho \rightarrow H \rtimes X^*$ is injective & a homomorphism

$$\begin{array}{l} x \mapsto (1, x) \\ g \mapsto (g, \varepsilon) \end{array}$$

□

Additional conditions

Theorem

Suppose $\mathcal{A}_{(H,X)}$ satisfies the serial processing condition, and both $\mathcal{A}_{(H,X)}$ and the dual automaton $\mathcal{A}_{(H,X)}^d$ are invertible. Then

$$\text{Mon}[H, X \mid xh = \theta_x(h)(x \cdot h) \ (h \in H, x \in X)]$$

is a Zappa-Szép monoid of H and X^* , and

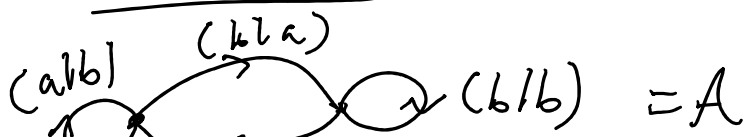
$$\text{Grp}[H, X \mid xh = \theta_x(h)(x \cdot h) \ (h \in H, x \in X)]$$

is a Zappa-Szép group of H and $F(X)$.

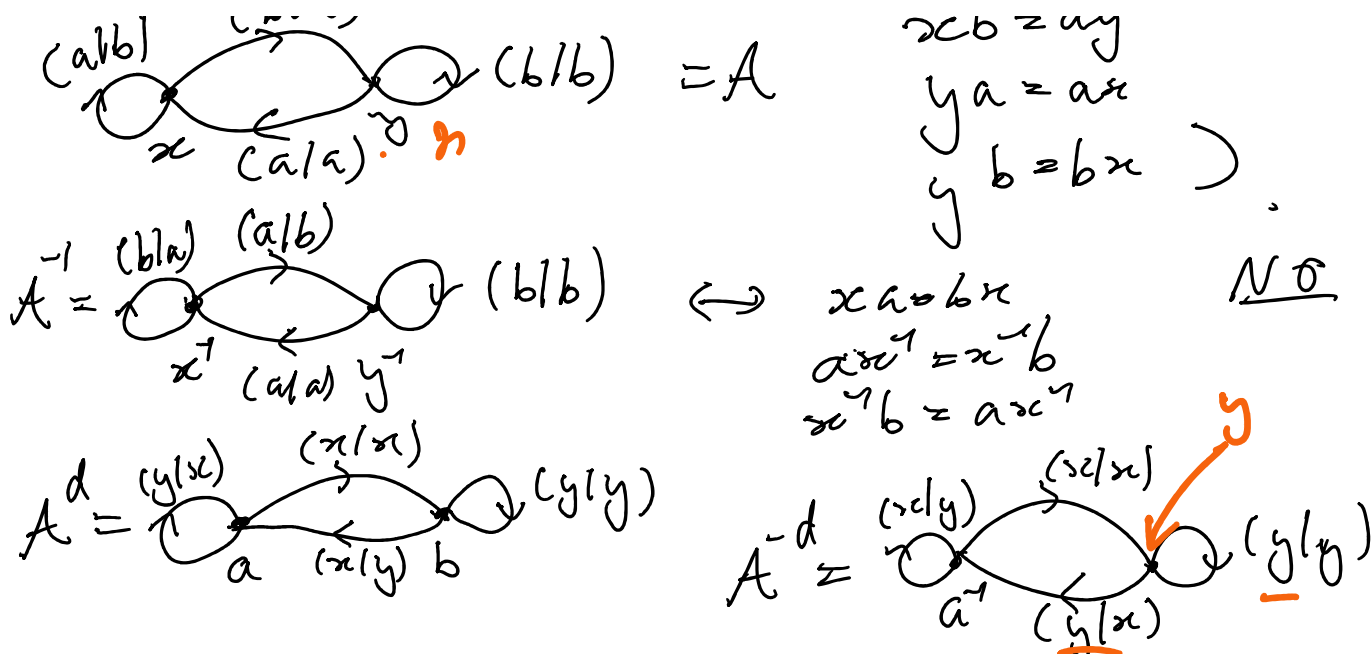
Dual and inverse automata

$\neq (a,b) \rtimes (x,y) \cdot \text{Grp} \langle a, b, x, y, z \mid$

non-example:



$$\begin{array}{l} xa = bx \\ xcb = ay \\ ya = az \end{array}$$



Application: Automorphism-induced HNN-extensions

Theorem

Fix G a group.

- Let $H \leq G$ be a subgroup.
- Let T be a transversal for G/H .
- Let $\phi \in \text{Aut}(G)$ be an automorphism.

Then the automorphism-induced HNN-extension

$$\langle G, t \mid tht^{-1} = \phi(h) \ (h \in H) \rangle$$

is isomorphic to the Zappa–Szép product

$$G \bowtie F(X) = \langle G, X \mid x_{\tau}g = \phi(g)x_{\tau\bar{g}} \ (\tau \in T, g \in G) \rangle.$$

Conclusions

- There is a link between Zappa–Szép products and automata for semigroups.
- This can be generalised and exploited to answer the question “when do a set of relators of the form $zx = \theta_z(x)(z \cdot x)$ give a Zappa–Szép product?”

Thank you for your attention!